

## Databases

### A little of mathematics

- We will be particularly interested in relational databases.
- Data are stored in tables.
- Why mathematics?  
Relational databases are inspired of relations - mathematics.

## Tables

- Set of rows (no duplicates)
- Each row describes a different entity.
- Each column states a particular fact about each entity.
  - Has an associated domain.

| Id   | Name | Address   | Status |
|------|------|-----------|--------|
| 1111 | John | 123 Main  | Fresh  |
| 2222 | Mary | 321 Oak   | Soph   |
| 1234 | Bob  | 444 Pine  | Soph   |
| 9999 | Joan | 777 Grand | Senior |

## Sets

- **Sets** are the basic data structures of mathematics.
- Intuitively, any “well-defined” collection of mathematical objects can be grouped together to form a single object – a *set*.
- Well-known examples of sets from mathematics are the set of integers ( $\mathbb{N}$ ), the set of rational numbers ( $\mathbb{Q}$ ), and the set of real numbers ( $\mathbb{R}$ ). In computer science one often deals with sets of strings (also called “formal languages”) or trees.

## Description of Sets

- A set with no elements is called an **empty set**. It is denoted  $\emptyset$  or  $\{\}$ . It is **unique**.

- The number of element of a set  $S$  is denoted  $|S|$ . We say also the **cardinal** of  $S$ .

- **Finite** sets can in principle be described by *listing* their elements.

That is, we write

$$\{x_1, \dots, x_n\}$$

to denote the set consisting of elements  $x_1, \dots, x_n$ .

- A more general mechanism for describing a set (*finite or infinite*) is to characterize via a *logical formula* a condition (property) its elements have to satisfy:

For every set  $S$  and formula  $P(x)$  there exists a set, denoted by

$$\{x \in S \mid P(x)\},$$

that consists of all elements of  $S$  for which  $P$  is true.

## Examples of Sets

- $\{\emptyset, 1, (4, 5), \text{"bonjour"}\}$

- The (finite) set of integers between  $-2$  and  $5$ :

$$\{n \in \mathbf{Z} \mid -2 < n < 5\}$$

- The (open) interval of real numbers between  $-2$  and  $5$ :

$$\{x \in \mathbf{R} \mid -2 < x < 5\}$$

- The (infinite) set of even integers:

$$\{n \in \mathbf{Z} \mid \exists k (n = 2k)\}$$

- From a general description it may not always be obvious what the elements of the set are:

–

$$\{(x, y, z) \in \mathbf{N} \times \mathbf{N} \times \mathbf{N} \mid (z = x + y)\}$$

–

$$\{(x, y, z) \in \mathbf{N} \times \mathbf{N} \times \mathbf{N} \mid \exists n \in \mathbf{N}, (n > 2 \wedge x^n + y^n = z^n)\}$$

## Set Theory

- The basic concepts of sets theory are **sets** and the **elements-relationship**.

- The symbol  $\in$  is commonly used to denote the **membership relation**, and one writes  $x \in A$  to denote the proposition  *$x$  is an element of  $A$*  (which may be true or false).

- Intuitively, sets are **unordered** collections of objects, where the **multiplicities** of elements *don't matter*.

- **Examples:**

$$\begin{aligned}\{1, 2\} &= \{2, 1\}? \\ \{1, 2\} &= \{1, 1, 2, 2, 2\}? \\ \{1, 2, 3\} &= \{1, 1, 1, 3\}?\end{aligned}$$

## Ordered Pairs and Tuples

- Sets are **unordered** collections of elements.

- **Pairs**, or more generally **tuples**, are **ordered** collections of elements.

- **Examples:**  $(1, 2)$  is a pair (a tuple of length 2).  $(1, 2, 4, 5)$  is a tuple of length 4.

- Tuples of different lengths are never equal.

- **Examples:**

$$\begin{aligned}(1, 2) &\neq (2, 1) \\ \{1, 2, 3\} &= \{1, 3, 2\} \\ (1, 2, 3) &\neq (1, 3, 2) \\ \{1, 2\} &= \{1, 2, 2\} \\ (1, 2) &\neq (1, 2, 2)\end{aligned}$$

## Subsets ( $\subseteq$ )

- **Definition:**

A set  $A$  is a **subset** of another set  $B$ , written  $A \subseteq B$ , if, and only if, every element of  $A$  is also an element of  $B$ .

- **Examples:**

$$\begin{aligned} \{1, 2\} &\subseteq \{1, 2, 3\}? \\ \{1, 1, 2, 2\} &\subseteq \{1, 2\}? \\ \{1\} &\subseteq \{2, 3, 5, 7\}? \end{aligned}$$

- The subset relation is often used to establish *equality* of sets, based on the following lemma.

**Lemma:** If  $A \subseteq B$  and  $B \subseteq A$ , then  $A = B$ .

## Proper subsets ( $\subset$ )

- **Definition:**

$A$  is a **proper subset** of  $B$ , written  $A \subset B$ , if  $A$  is a subset of  $B$ , but not equal to  $B$ .

- **Examples:**

$$\{1, 2\} \subset \{1, 2, 3\}?$$

$$\{1, 2\} \subset \{1, 1, 2, 2\}?$$

## Membership and subset relations

- Be careful about the distinction between the element relation and the subset relation.

- **Examples:**

$$\begin{aligned} 2 &\in \{1, 2, 3\}? \\ \{2\} &\in \{1, 2, 3\}? \\ 2 &\subseteq \{1, 2, 3\}? \\ \{2\} &\subseteq \{1, 2, 3\}? \\ \{2\} &\subseteq \{\{1\}, \{2\}\}? \\ \{2\} &\in \{\{1\}, \{2\}\}? \end{aligned}$$

## Property of the Empty Set

- **Theorem:**

If  $\emptyset$  is an empty set, then  $\emptyset \subseteq A$ , for all sets  $A$ .

## More Set Operations

### Cartesian Products

- Pairs and tuples provide us with a way of constructing new sets from given ones.

- **Definition:**

If  $A$  and  $B$  are sets, then there exists a set  $A \times B$  (read "A cross B"), called the *Cartesian product* of  $A$  and  $B$ , that consists of all ordered pairs  $(a, b)$ , where  $a \in A$  and  $b \in B$ .

- Symbolically,

$$A \times B = \{(a, b) \mid a \in A \wedge b \in B\}.$$

- For example, if  $A = \{1, 2\}$  and  $B = \{4, 5\}$ , then

$$A \times B = \{(1, 4), (1, 5), (2, 4), (2, 5)\}.$$

### Properties of the Empty Set

- **Theorems:**

$$\begin{aligned} A \cup \emptyset &= A \\ A \cap \emptyset &= \emptyset \end{aligned}$$

- Other operations for constructing sets include
  - *set union* ( $\cup$ )
  - *set intersection* ( $\cap$ )
  - *relative complementation* (or *set difference*) ( $-$ )
  - *complementation* ( $^c$ )

They are defined as follows.

- Let  $A$  and  $B$  be subsets of some set  $U$ . We define:

$$\begin{aligned} A \cup B &= \{x \in U \mid x \in A \vee x \in B\} \\ A \cap B &= \{x \in U \mid x \in A \wedge x \in B\} \\ B - A &= \{x \in U \mid x \in B \wedge x \notin A\} \\ A^c &= \{x \in U \mid x \notin A\} \end{aligned}$$

Note that set difference can also be defined as follows:

$$A - B = A \cap B^c.$$

- For example, let

$R$  be the set of real numbers,  
 $A$  the set  $\{x \in \mathbf{R} \mid -1 < x \leq 0\}$ ,  
 $B$  the set  $\{x \in \mathbf{R} \mid 0 \leq x < 1\}$ .

What are  $A \cup B$ ,  $A \cap B$ ,  $B - A$ , and  $A^c$ ?

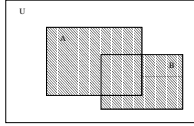
### Set Identities

- 1. Set union and intersection are commutative.
- 2. Set union and intersection are associative.
- 3. Distributivity:  $A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$
- 4. Double complement:  $(A^c)^c = A$ .
- 5. Idempotency:  $A \cap A = A \cup A = A$ .
- 6. De Morgan's Laws:  
 $(A \cup B)^c = A^c \cap B^c$  and  $(A \cap B)^c = A^c \cup B^c$ .
- 7. Absorption:  $A \cup (A \cap B) = A$  and  $A \cap (A \cup B) = A$ .

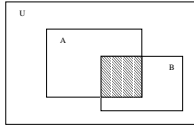
## Venn Diagrams

- Sets can often be conveniently represented by **Venn diagrams**.

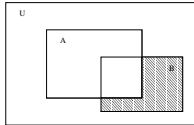
- The union  $A \cup B$  of  $A$  and  $B$  is represented by:



- The intersection  $A \cap B$  is represented by:



- The set difference  $B - A$  is represented by:



## Powersets ( $\mathcal{P}$ )

- **Powerset Axiom:**

If  $A$  is a set, then there exists a set, called the **powerset** of  $A$  and denoted by the symbol  $\mathcal{P}(A)$ , whose elements are exactly all the subsets of  $A$ .

- **Example:**

If  $A$  is the set  $\{1, 2, 3\}$ , then

$$\mathcal{P}(A) = \{ \emptyset, \{1\}, \{2\}, \{3\}, \{1, 2\}, \{1, 3\}, \{2, 3\}, \{1, 2, 3\} \}$$

Do we have  $1 \in \mathcal{P}(A)$ , or  $2 \in \mathcal{P}(A)$ , or  $3 \in \mathcal{P}(A)$ ?

No, because  $1 \neq \{1\}$ , etc.

- If  $A$  has  $n$  elements, how many elements are there in its powerset?

Answer:  $2^n$ . Why?

## Disjoint Sets

- Two sets  $A$  and  $B$  are said to be **disjoint** if they have no elements in common, i.e.,  $A \cap B = \emptyset$ .

- **Examples:**

Is  $\{\emptyset, \{\emptyset\}\} \cap \{\emptyset\} = \emptyset$ ?

No,  $\{\emptyset, \{\emptyset\}\} \cap \{\emptyset\} = \{\emptyset\}$ .

Is  $\{\emptyset, \{\emptyset\}\} \cap \emptyset = \emptyset$ ?

Yes, because  $A \cap \emptyset = \emptyset$ .

- A **partition** of a set  $A$  is a collection of pairwise disjoint sets  $A_1, \dots, A_n$ , such that

$$A = A_1 \cup A_2 \cup \dots \cup A_n.$$

- For example, at the end of the semester I will partition the class into subsets with grades of  $A$ ,  $A-$ , etc. It will be a partition, since each student gets one, and only one, grade.

## Relations

- **Relations** use ordered tuples to represent relationships among objects.

- **Examples:**

– “ $x$  is a parent of  $y$ ” –  $(\text{Morris}, \text{Steve}), (\text{Ria}, \text{Steve})$

– “ $x$  is a number less than  $y$ ” –  $(3, 42), (42, 43)$

– “Student number  $x$  is named  $y$  and majors in  $z$ ” –  $(124324443, \text{Mary}, \text{CSE}), (563565426, \text{Mary}, \text{PSY})$

– “ $x$  is an even number” ... (2)

- Essentially, a relation is the set of assignments which makes a predicate true.

- **Examples:**

–  $\text{IsParent} = \{(\text{Morris}, \text{Steve}), (\text{Ria}, \text{Steve})\}$

–  $\text{LessThan} = \{(3, 42), (42, 43)\}$

–  $\text{MajorIn} = \{(124324443, \text{Mary}, \text{CSE}), (563565426, \text{Mary}, \text{PSY})\}$

–  $\text{IsEven} = \{n \mid n = 2k\}$

# Presenting Binary Relations

## Binary Relations

- Binary relations have two blanks, relating two objects.
- More formally, suppose  $A$  and  $B$  are sets.  
A **binary relation** from  $A$  to  $B$  is a set  $R \subseteq A \times B$ .
- Thus  $R$  is a set of ordered pairs  $(a, b)$  where  $a \in A$  and  $b \in B$ .
- **Notation:** If  $(a, b) \in R$  then we sometimes write  $aRb$ .
- **Example:**  
 $A = \{2, 6, 7\}$ ,  $B = \{1, 2, 5\}$ .  
 $R_1$  is "x in  $A$  is an integer multiple of y in  $B$ ."  
so  $R_1 = \{(2, 1), (2, 2), (6, 1), (6, 2), (7, 1)\}$

- Binary relations are particularly useful because they have two kinds of compact visual representation, **tables** and **graphs**.

- **Tables:**

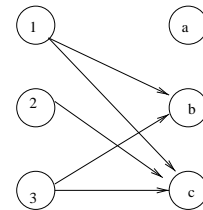
|   |   |   |   |
|---|---|---|---|
| R | a | b | c |
| 1 |   | * | * |
| 2 |   |   | * |
| 3 |   | * | * |

or

|   |   |
|---|---|
| R | b |
| 1 | b |
| 1 | c |
| 2 | c |
| 3 | b |
| 3 | c |

- **Graphs** are composed of **vertices** or **nodes** connected by **edges** or **arcs**.

There is an arc from  $a$  to  $b$  iff  $(a, b) \in R$

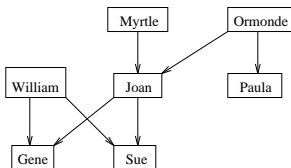


## The Parent-Of Relation

- The parent of relations, "x is a parent of y", is a binary relation between pairs of people.
- **Table:**

| R       | Gene | Joan | William | Sue | Myrtle | Ormonde | Paula |
|---------|------|------|---------|-----|--------|---------|-------|
| Gene    |      |      |         |     |        |         |       |
| Joan    | *    |      |         |     |        |         |       |
| William | *    |      |         |     |        |         |       |
| Sue     |      |      |         |     |        |         |       |
| Myrtle  |      | *    |         |     |        |         |       |
| Orm     |      | *    |         |     |        |         |       |
| Paula   |      |      |         |     |        |         | *     |

- **Graph:**



- Which representation is better for testing whether the pair  $(x, y)$  is in the relation?
- Which representation is better for capturing the overall structure?

## General ( $n$ -ary) Relations

- Suppose  $A_1, A_2, \dots, A_n$  are sets. A relation of  $A_1, A_2, \dots, A_n$  is a set  $R \subseteq A_1 \times A_2 \times \dots \times A_n$ .

- Thus  $R$  is a set of ordered  $n$ -tuples  $(a_1, a_2, \dots, a_n)$  where  $a_i \in A_i$ .

- **Example:**

$A_1 = N$ ,  $A_2 = names$ ,  $A_3 = majors$

"Student number x is named y and majors in z"

$(124324443, Mary, CSE)$ ,  $(563565426, Mary, PSY)$  are tuples of the relation.

- Such structures are modeled by **hypergraphs**, a graph structure where each "edge" represents a subset of more than two vertices.

# Relational Databases

## Overview

- The most important commercial database systems today employ the **relational** model, meaning that the data is stored as tables of tuples, i.e. relations.

A relation is a mathematical entity corresponding to a table:

- row - tuple
- column attribute.

- A Shakespearian **killed** relation would be:

| Killer  | Victim   |
|---------|----------|
| Brutus  | Caesar   |
| Hamlet  | Laertes  |
| Hamlet  | Polonius |
| Laertes | Hamlet   |
| Brutus  | Brutus   |
| Cassius | Caesar   |

- **Requests** for information from the database is made in a query language like **SQL** which is based on the notations of set theory and the predicate calculus.

- **Example 1: Who killed Caesar?**

- **Example 2: Who was both a killer and a victim?**

In SQL:

```
(SELECT Killer from Killed) INTERSECT (SELECT Victim from Killed)
```

- Much of the power of relational databases comes from the fact that we can **combine different relations**.
- For example, suppose we also have a **died-by** relation:

| Victim   | Method  |
|----------|---------|
| Caesar   | Daggers |
| Hamlet   | Sword   |
| Laertes  | Sword   |
| Polonius | Sword   |
| Brutus   | Sword   |

We can combine the two tables with a **join** operation, which the tables based on **common fields**. For example, the join of **killed** and **died-by** is:

| Killer  | Victim   | Method  |
|---------|----------|---------|
| Brutus  | Caesar   | Daggers |
| Hamlet  | Laertes  | Sword   |
| Hamlet  | Polonius | Sword   |
| Laertes | Hamlet   | Sword   |
| Brutus  | Brutus   | Sword   |
| Cassius | Caesar   | Daggers |

In SQL:

```
SELECT Killer from Killed where victim='Caesar'
```

This reads "select from relation 'killed' all tuples where the victim was Caesar, and report only the killer field from each.

- **Example 3: Which killers used daggers?**

In SQL:

```
SELECT Killer FROM Killed, Died_by  
WHERE Killed.victim=died_by.victim  
AND Method='Daggers'
```

- Note that this database design assumes that each victim can only be killed by one weapon (sorry, Rasputin).